

# Implementing a Toffoli Gate in an N-Qubit Linear Nearest Neighbor Array

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**Abstract.** Most proposals for quantum computers are based on linear nearest neighbor (LNN) arrangements where qubits only interact with neighbors. A Toffoli gate is a three-qubit operation, which is used in many quantum applications such as error correction, and algorithms like Shor’s factorization. Typically, to implement a Toffoli gate in an LNN architecture, additional operations called swap gates are required to bring the qubits adjacent to each other. We propose a new method to implement a Toffoli gate in LNN arrays without using swap gates. As such, compared to other circuits, the quantum cost of our circuit is much lower.

## 1. Introduction

In quantum computing, a qubit (or quantum bit) is represented by the state  $\alpha|0\rangle + \beta|1\rangle$  where  $|\alpha|^2 + |\beta|^2 = 1$ . A Toffoli gate is a three-qubit gate, which has two controls and one target qubit. When the two controls are each in the  $|1\rangle$  state, the state of the target is flipped. Currently, to implement a Toffoli gate in a three-qubit system 11 elementary gates are required (single qubit gates and controlled-NOT gates) [1]. As the number of qubits increases, the total number of elementary gates increases. This is because most physical arrangements of qubits are along an LNN array where a qubit only interacts with its neighbors. As such, to interact two non-NN qubits, additional gate operations are required to bring them adjacent to each other. Here, we show methods for implementing a three-qubit Toffoli gate efficiently in an N-qubit system.

## 2. Toffoli gate in a three qubit system

Fig. 1 shows a Toffoli gate for a three-qubit system with qubit 3 as the target. Here, we implement the gate by sandwiching a  $C^2(-I)$  gate between two controlled-Hadamard ( $C^1(H)$ ) gates. A  $C^2(-I)$  gate has two controls and 1 target [2]. A  $\pi$  phase is picked up by the target only when the two controls are each in the  $|1\rangle$  state. Overall, a  $C^2(-I)$  gate is equivalent to a

controlled-Z ( $C^1(Z)$ ) gate, which is a symmetric gate (either qubit can be the target or the control). The  $C^1(Z)$  and  $C^1(H)$  gates implement a Z and an Hadamard (H) gate, respectively, on the target qubit when the control qubit is in the  $|1\rangle$  state. The Z and H gates are defined as:

$$Z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix} ; \quad H = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} \quad (1)$$

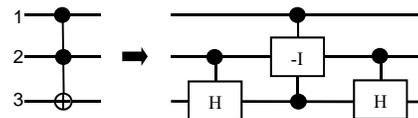


Fig.1. Implementation of a Toffoli gate with qubit 3 as the target, using 3 elementary gates [2].

## 3. Toffoli gate in a four qubit system

Fig. 2 shows our method for implementing a Toffoli gate in a four qubit system where two of participants are not adjacent. Hereafter, we will refer to the control and target qubits as participants. The Toffoli gate is realized by sandwiching a  $C^2(Z)$  gate between two H gates applied on the target qubit. For realizing the  $C^2(Z)$  gate, two Toffoli gates between qubits 1, 2 and 3 with qubit 2 as the target and two  $C^1(Z)$  gates between qubits 2 and 4, are implemented. The two  $C^1(Z)$  gates between qubits 2 and 4 are replaced by two  $C^2(-I)$  gates (Fig. 2).

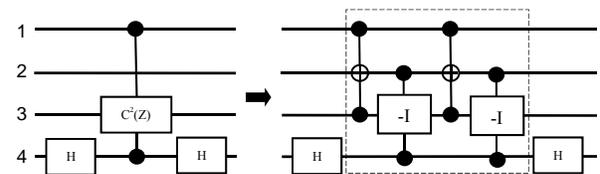


Fig. 2. Implementation of a Toffoli gate in a four qubit system using qubit 2 as the dummy.

### 4. Toffoli gate in an $N$ qubit system

Fig. 3 shows a Toffoli gate with three participants in an  $N$  qubit system, where qubits  $p$ ,  $q$  and  $r$  are the participants. Here and throughout,  $p$  is the 1<sup>st</sup> qubit,  $q$  is in the middle, and  $r$  is the  $N^{\text{th}}$  qubit. As before, to perform a Toffoli gate, a  $C^2(Z)$  gate is realized on these three participants which is sandwiched between two H gates applied on the target qubit. In our method, we divide qubits into two major groups of “ $m$ ” and “ $(N-m)+1$ ” qubits. If  $N$  is odd, we choose “ $m=(N+1)/2$ ”, and if  $N$  is even, we choose “ $m=(N+2)/2$ ”. In either case, if  $m=q$ , we change  $m=q+1$ . The dummy qubit is selected on  $m$ 's position (a dummy qubit is one that is used as a target that is finally restored to its initial state). Depending on whether participant  $q$  is in group “ $m$ ” (case 1) or in group “ $(N-m)+1$ ” (case 2), each  $C^2(Z)$  gate is decomposed into four gates. In case 1, the  $C^2(Z)$  gate is replaced by two Toffoli gates between  $p$ ,  $q$  and  $m$ , with  $m$  as the target, and two  $C^1(Z)$  gates between  $m$  and  $r$ . In case 2, the  $C^2(Z)$  gate is substituted by two Toffoli gates between  $r$ ,  $q$  and  $m$  ( $m$  as the target), and two  $C^1(Z)$  gates between  $m$  and  $p$ .

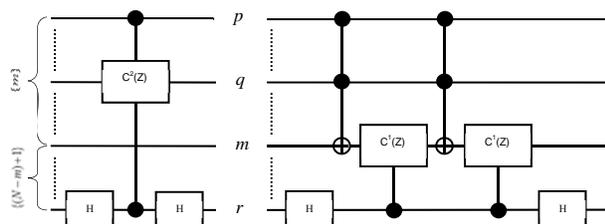


Fig. 3. Implementation of a Toffoli gate in an  $N$  qubit system (case 1) where participant  $q$  is in group “ $m$ ”.

In Fig. 3, the two Toffoli implemented between qubits  $p$ ,  $q$  and  $m$  (group “ $m$ ”) need not be between NN qubits. If the Toffoli gates are not of the form of Fig. 1, further decomposition into sub-groups may be needed. This process of decomposition is carried out until all Toffoli gate operations are of the form of Fig. 1. Likewise, in each group/subgroup, if the  $C^1(Z)$  gates are such that the separation between the control and target qubits is more than one qubit, each gate will have to be broken down into subgroups. In Ref. [2], the author shows how to implement a  $C^1(Z)$  gate between non-NN qubits. For  $P \geq 7$ , where  $P$  is the number of qubits, the total gate count for the  $C^1(Z)$  gate was shown to be [2]:

$$G_p = (20 + (P - 7) \times 6); P \geq 7 \quad (2)$$

Note that, here  $P = (N-m)+1$ .

Following the procedure of decomposing the  $N$  qubit system into groups and subgroups, the total gate count for the Toffoli gate was calculated for different  $N$ . Fig. 4 shows the plot (blue line). From the plot, a second order polynomial trend line was calculated (black line), which gave an equation for the gate count,  $G_N$  (total number of elementary gates) as a function of  $N$ :

$$G_N = 0.2519N^2 + 6.6723N - 36.973 \quad (3)$$

As a measure of the closeness of this equation to the actual gate count, for  $N = 40$ , Eq. (3) gives  $G_N = 633$ , while the exact value of  $G_N$  using our method is 626.

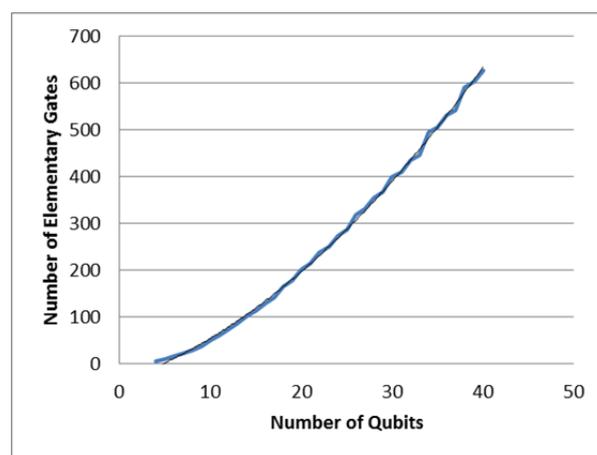


Fig. 4. Number of qubits versus number of elementary gates.

Observing Eq. (2),  $G_N$  is a quadratic function of  $N$ . The gate count can be further reduced by making it a linear function of  $N$ . This can be done by combining our method with swap operations. It was found that for  $N \geq 7$ :

$$G_N = (22 + (N - 7) \times 6); N \geq 7 \quad (4)$$

### 5. Conclusions

In this paper, we introduced a new method to implement a Toffoli gate in an LNN array efficiently, by reducing the number of elementary gates. A quadratic as well a linear variation of the total gate count with the number of qubits has been presented.

[1] Mehdi Saeedi, Robert Wille, Rolf Drechsler, Quantum Information Processing June 2011, Volume 10, Issue 3, pp 355-377  
 [2] Preethika Kumar, Quantum Information Processing April 2013, Volume 12, Issue 4, pp 1737-1757